### CS152: Computer Systems Architecture Moore's Law

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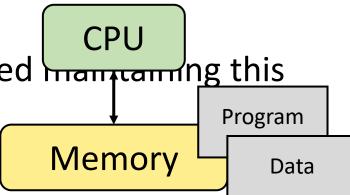


Gordon Moore (Photo 1978) 1929 – 2023



#### Conventional performance scaling

- ☐ Traditional model of a computer is simple
  - Single, in-order flow of instructions on a processor
  - Simple, in-order memory model
- Large part of computer architecture research involved mangaming this abstraction while improving performance
  - Transparent caches, Transparent superscalar scheduling,
  - Same software runs faster tomorrow
  - (Slow software becomes acceptable tomorrow)
- ☐ Driven largely by continuing march of Moore's law



#### Moore's Law

- ☐ What exactly does it mean?
- ☐ What is it that is scaling?

#### Moore's Law

☐ Typically cast as:

"Performance doubles every X months"

☐ Actually closer to:

"Number of transistors per unit cost doubles every X months"

#### Moore's Law

The complexity for minimum component costs has increased at a rate of roughly a factor of two per year.

[...]

Over the longer term, the rate of increase is a bit more uncertain, although there is no reason to believe it will not remain nearly constant for at least 10 years.

-- Gordon Moore, Electronics, 1965

Why is Moore's Law conflated with processor performance?

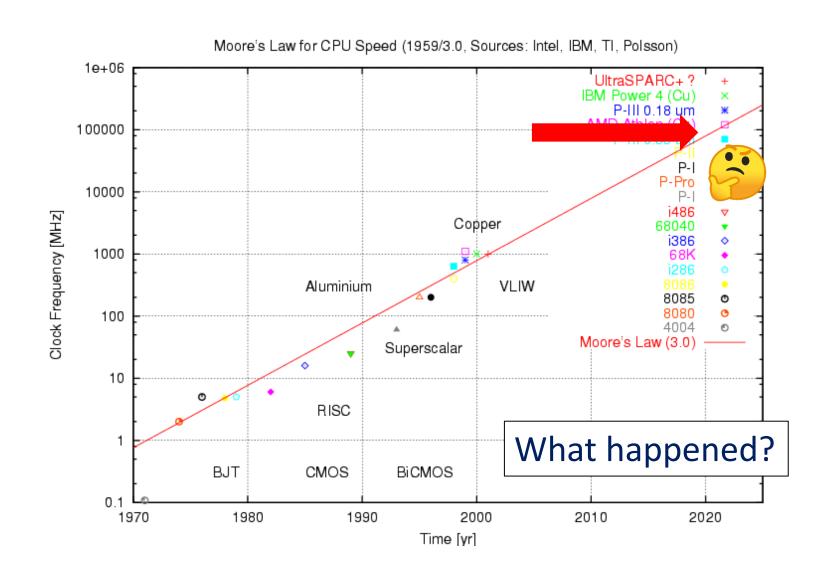
## Dennard Scaling: From Moore's Law to performance

- ☐ "Power density stays constant as transistors get smaller"
  - Robert H. Dennard, 1974

- ☐ Intuitively:
  - $\circ$  Smaller transistors  $\rightarrow$  shorter propagation delay  $\rightarrow$  faster frequency
  - Smaller transistors → smaller capacitance → lower voltage
  - Power  $\propto$  Capacitance  $\times$  Voltage<sup>2</sup>  $\times$  Frequency

Moore's law → Faster performance @ Constant power!

### Single-core performance scaling projection



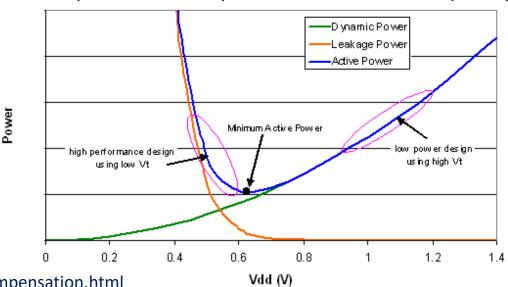
# (Slightly) more accurate processor power consumption

Gate-oxide Stopped scaling stopped scaling due to leakage  $Power = (ActiveTransistors \times Capacitance \times Voltage^2 \times Frequency)$ 

Dynamic power

Total power consumption with constant frequency





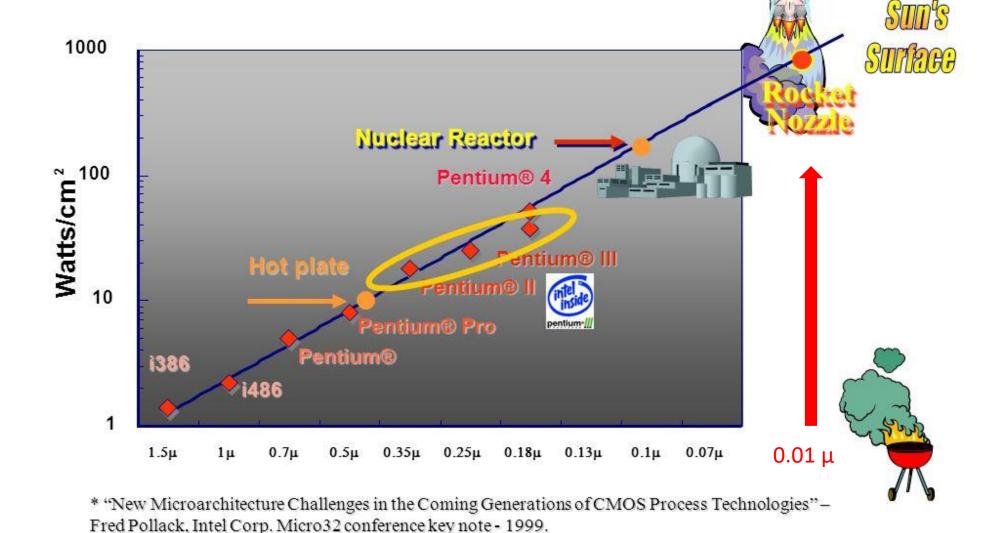
#### End of Dennard Scaling

- ☐ Even with smaller transistors, we cannot continue reducing power
  - O What do we do now?

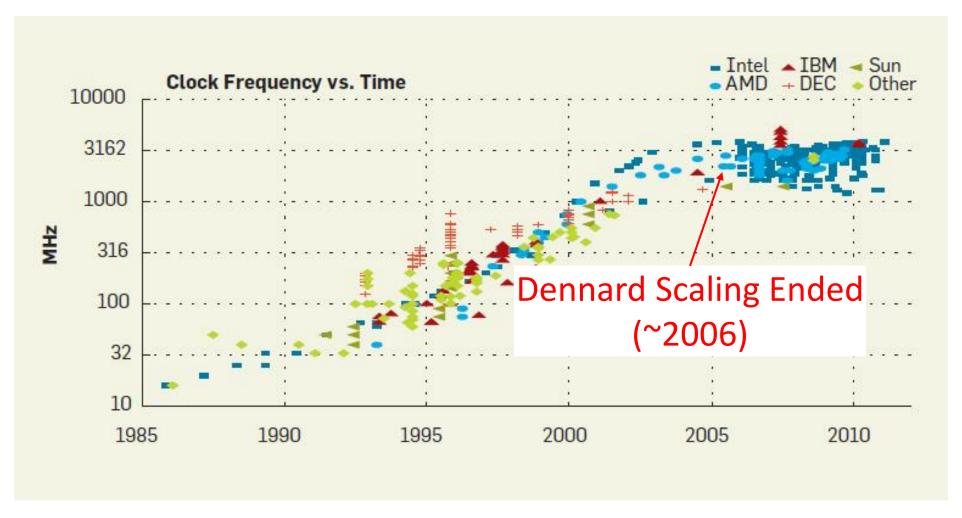
- Option 1: Continue scaling frequency at increased power budget
  - Chip quickly become too hot to cool!
  - Thermal runaway:

Hotter chip  $\rightarrow$  increased resistance  $\rightarrow$  hotter chip  $\rightarrow$  ...

Option 1: Continue scaling frequency at increased power budget

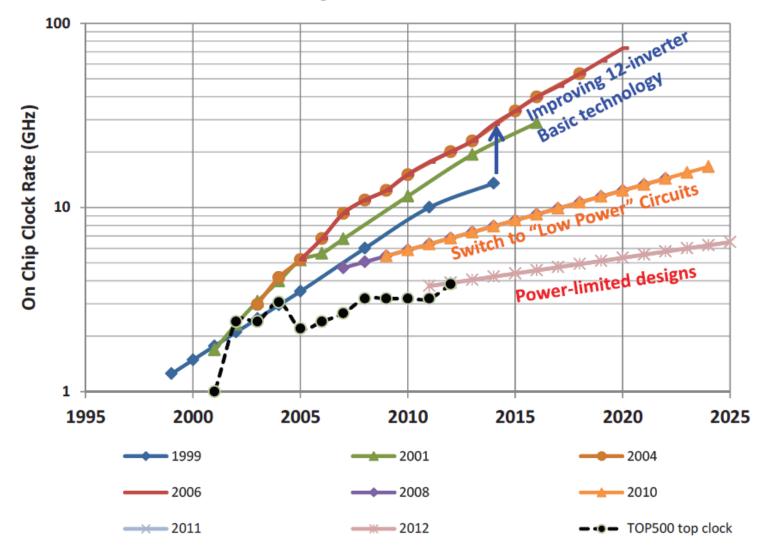


### Option 2: Stop frequency scaling



Danowitz et.al., "CPU DB: Recording Microprocessor History," Communications of the ACM, 2012

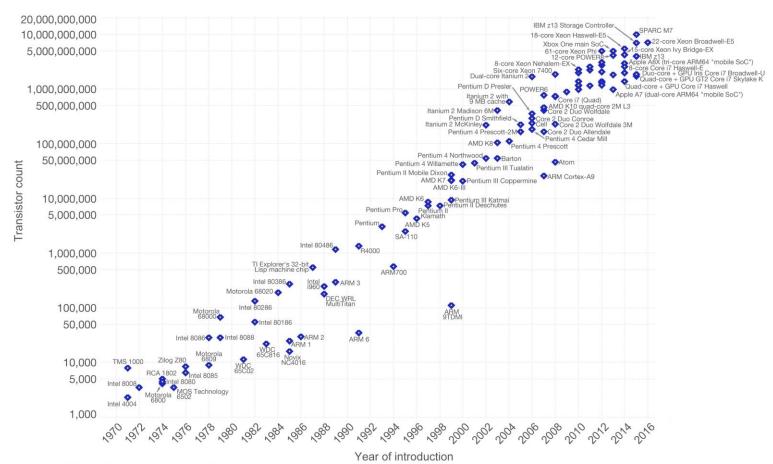
#### Looking back: change of predictions



### But Moore's Law continues beyond 2006

#### Moore's Law – The number of transistors on integrated circuit chips (1971-2016)



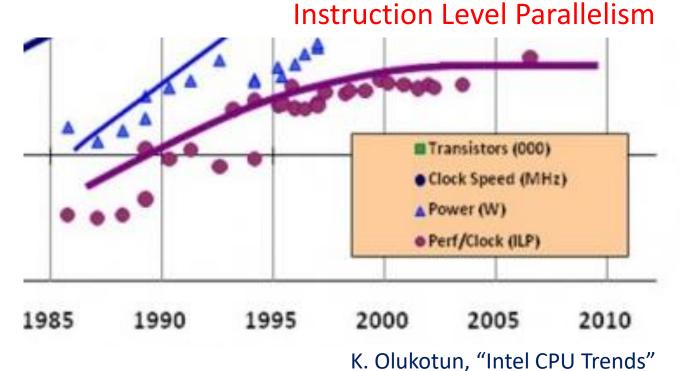


### State of things at this point (2006)

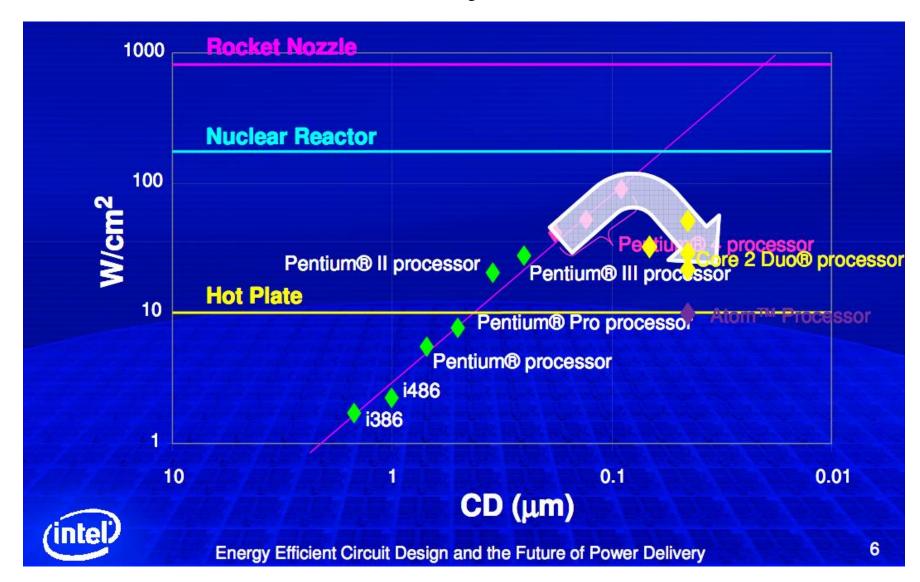
- ☐ Single-thread performance scaling ended
  - Frequency scaling ended (Dennard Scaling)
  - Instruction-level parallelism scaling stalled ... also around 2005

#### ☐ Moore's law continues

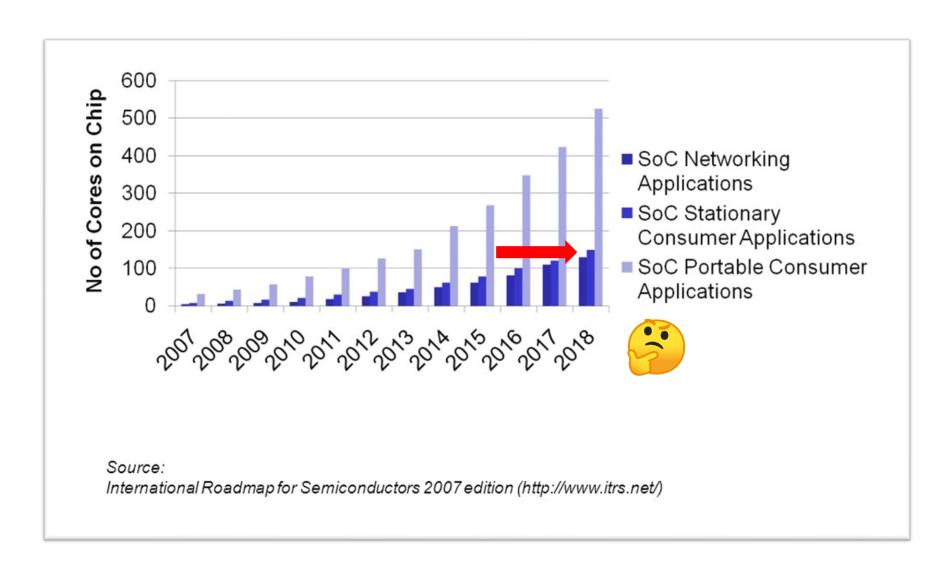
- Double transistors every two years
- Owner with White White or with them?



#### Crisis averted with manycores?



#### Crisis averted with manycores?



#### What happened?

```
Can't keep going up

Gate-oxide Stopped scaling Stopped scaling stopped scaling due to leakage due to thermal

(ActiveTransistors \times Capacitance \times Voltage^2 \times Frequency)

Dynamic power
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+ (Voltage × LeakageCurrent)

"Utilization Wall"

Static power

Regardless of Moore's Law, a limited amount of gates can be active at a given time

#### Where To, From Here?

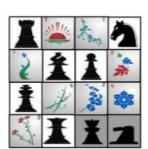
- ☐ The number of active transistors at a given time is limited
  - We won't get much performance improvements even if Moore's law continues
  - We need to make the best use of those active transistors!

#### Where To, From Here?

- ☐ Potential Solution 1: The software solution
  - Write efficient software to make the efficient use of hardware resources
  - No longer depend entirely on hardware performance scaling
  - "Performance engineering" software, using hardware knowledge
- ☐ Solution 2: The specialized architectural solution
  - Chip space is now cheap, but power is expensive
  - Stop depending on more complex general-purpose cores
  - Use space to build heterogeneous systems,
     with compute engines well-suited for each application







## The Bottom Line: Architecture is No Longer Transparent

- Optimized software requires architecture knowledge
- ☐ Special-purpose "accelerators" (GPU, FPGA, ...) programmed explicitly
- Even general-purpose processors implement specialized instructions
  - Single-Instruction Multiple Data (SIMD) instructions such as AVX
  - Special-purpose instructions sets such as AES-NI